

Types

COMP360

Some slides are from

“Programming Language Paradigms” by Michael Scott

“It’s hardware that makes a machine fast. It’s software that makes a fast machine slow.”

Craig Bruce

Bring your laptop

- On Wednesday we will be discussing Drag & Drop programming
- It would be useful to have you try some examples during the lecture
- read
<http://appinventor.mit.edu/explore/ai2/tutorials.html>

Remaining Schedule

Monday, April 17 Types read sections 7.1 & 7.2	Wednesday, April 19 Drag & drop programming	Friday, April 21 Drag & drop programming
Monday, April 24 Concurrent Programming read chapter 13	Wednesday, April 26 Concurrent Programming	Friday, April 28 Concurrent Programming
Monday, May 1 Exam 3	Wednesday, May 3 final review	
Tuesday, May 9 Final Exam 10:30am – 12:30pm		

Data Format

- Different types of data are stored differently
- Some operations are not allowed on some data types

What is $5 * \text{“aardvark”}$?

- Some operations are implemented differently for different data types

Data Types

- What are types good for?
 - implicit context
 - checking - make sure that certain meaningless operations do not occur
 - type checking cannot prevent all meaningless operations
 - It catches enough of them to be useful
- Polymorphism results when the compiler finds that it doesn't need to know certain things

Old Languages Without Types

- In the pre-Cambrian days of programming, C didn't really care about types
- Functions returning an integer or pointer with the same size as an **int** are not required to have a declaration
- Parameters did not have to be declared

Modern Example

- Consider this C function written in a modern style

```
double new_style( int cat, double *dog); /* prototype */
```

```
double new_style( int cat, double *dog) {  
    return *dog + (double)cat;  
}
```


Old C Example

- Consider this C function written in an ancient style

```
double old_style( );      /* prototype */
```

```
double new_style( cat, dog )  
    int cat, double *dog  
{  
    return *dog + (double)cat;  
}
```

Strong Typing

- **Strong Typing** has become a popular buzz-word
 - like *structured programming*
 - informally, it means that the language prevents you from applying an operation to data on which it is not appropriate
- **Static Typing** means strongly typed and that the compiler can do all the checking at compile time

Example Type Systems

- Common Lisp is strongly typed, but not statically typed
- Ada is statically typed
- Pascal is almost statically typed
- Java is strongly typed, with a non-trivial mix of things that can be checked statically and things that have to be checked dynamically
- C has become more strongly typed with each new version, though loopholes still remain

Common Hardware Types

- Most CPUs provide instructions to handle certain data types
 - integer
 - long
 - short
 - char
 - float
 - double
 - pointers

Software Created Types

- Arrays
- Strings
- Classes
- Records
- Sets
- Files
- Enumeration

Coercion

- When an expression of one type is used in a context where a different type is expected, one normally gets a type error
- But what about

```
int cat;   double dog, goat;
```

```
...
```

```
dog = cat + goat;
```



Which Java statement will work?

```
double dog, goat;    int cat;
```

- A. `cat = dog + goat;`
- B. `dog = cat + goat;`
- C. `cat = (int)dog + goat;`
- D. All the above
- E. None of the above

Coercing Expressions

- Many languages allow things like this, and COERCE an expression to be of the proper type
- Coercion can be based just on types of operands, or can take into account expected type from surrounding context as well
- Fortran has lots of coercion, all based on operand type

Simple Coercion

- C has lots of coercion, too, but with simpler rules:
 - all **floats** in expressions become **doubles**
 - **short**, **int**, and **char** become **int** in expressions
 - if necessary, precision is removed when assigning into LHS

No Static Types

- Some languages, such as JavaScript, do not statically type the variables
- A variable can hold any data given to it

```
var rabbit;  
rabbit = 7;  
rabbit = "hare";
```

- Data types are checked at run time

Generics

- Generic classes are particularly valuable for creating *containers*: data abstractions that hold a collection of objects
- Generics allow type determination without casting
- Generic methods are needed in generic classes

Generics Example

- Casting is required without generics

```
ArrayList cow = new ArrayList();  
String bull = (String) cow.get();
```

```
ArrayList<String> goat = new ArrayList<>();  
String bull = goat.get();
```

Strings

Usual string operations include:

- Concatenation
- Substring
- Comparison
- Length
- Assignment
- *Pattern matching*

String Design Consideration

- Character arrays
- Length
 - Static
 - Dynamic
 - Limited length dynamic
- Immutable

String Implementation

- Length determination
 - Special termination character
 - Property
 - Stored separately from the string data
 - Stored with the string data
- How do you handle changes in string length?

Immutable String

- Java uses immutable strings
- You cannot change a string although you can replace one

```
String cat = "meow";
```

```
String dog = cat;
```

```
cat = cat + " purr";
```

```
System.out.println( dog );    // displays meow
```


Array Considerations

- Specification of index range
- Dynamic or Static size
- Number of dimensions
- Range checking
- Initialization
- Slices

Array Organization

- Vectors are almost always stored as consecutive addresses.
- 2 dimensional arrays can be:
 - Consecutive addresses
 - Array of arrays
- 3+ dimensional arrays

2D Array Allocation

```
int a[3][2];
```

0,0	0,1
1,0	1,1
2,0	2,1

Allocation in memory

0,0	0,1	1,0	1,1	2,0	2,1
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2D Array Index Calculation

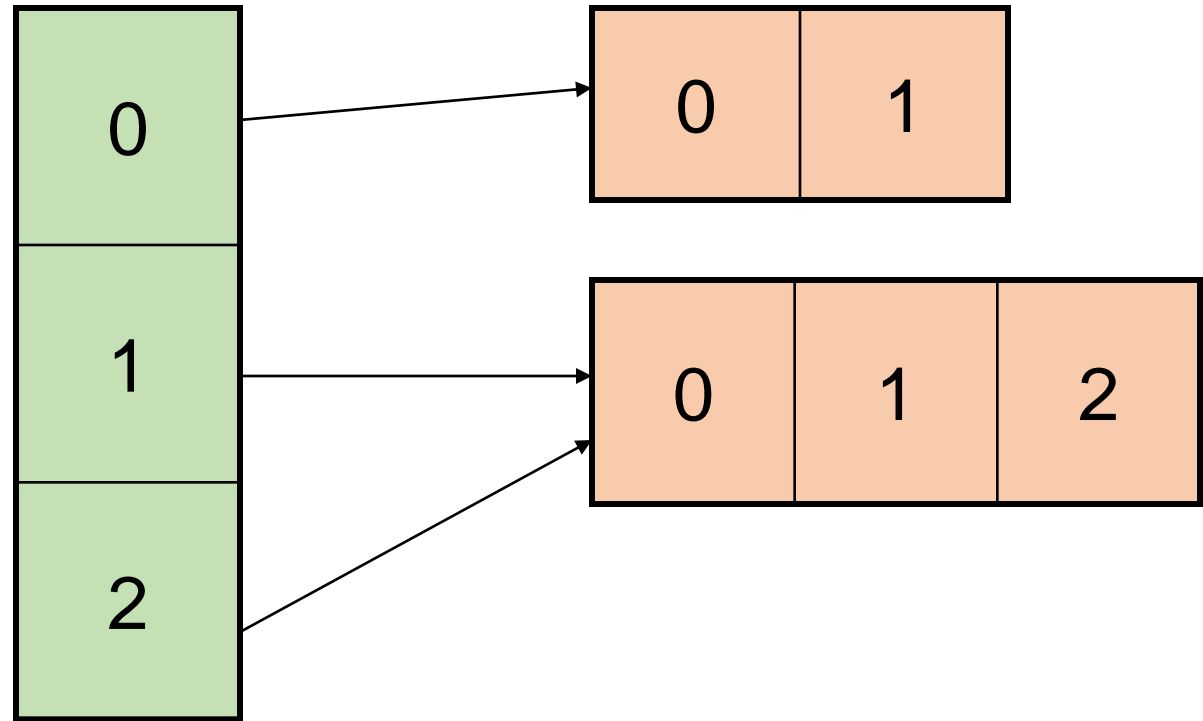
Location of element $i,j =$

Base address +

object size * num col * $i +$

object size * $j;$

Java Arrays of Arrays



Array Slices

- Some languages allow you to operate with an arbitrary portion of an array
- C allows you to pass a row of a 2D array to a function as a vector, but you can't pass a column as a vector. Fortran 90 allows this

Orthogonality

- **Orthogonality** is a useful goal in the design of a language, particularly its type system
- A collection of features is orthogonal if there are no restrictions on the ways in which the features can be combined
- Orthogonality is nice primarily because it makes a language easy to understand, easy to use, and easy to reason about

Non-Orthogonality in C

- Structures (but not arrays) may be returned from a function
- An array can be returned if it is inside a structure
- A member of a structure can be any data type (except void), or the structure of the same type
- An array element can be any data type (except void)
- Everything is passed by value (except arrays)

Type Inference

- Haskell is strongly typed
- While you can, you are not required to specify the type of variables
- Haskell determines the type of the variables based on the operations done on them

Type Checking

A type system has rules for

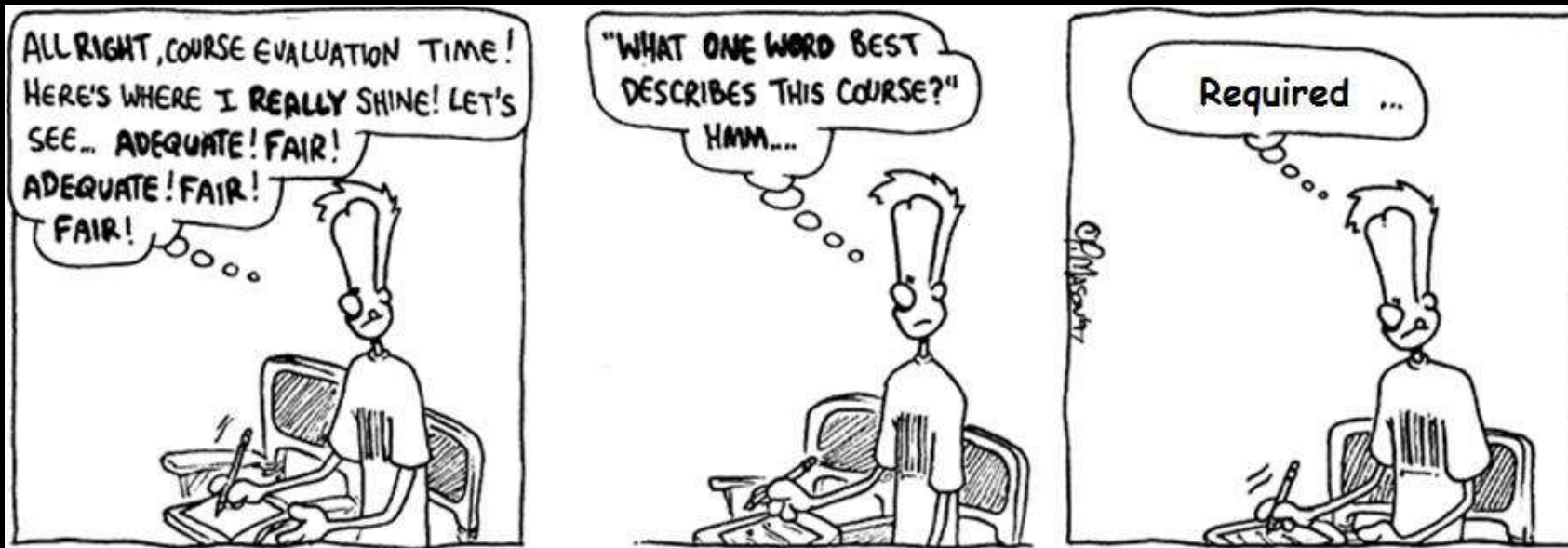
- **type equivalence** – when are the types of two values the same?
- **type compatibility** – when can a value of type A be used in a context that expects type B?
- **type inference** – what is the type of an expression, given the types of the operands?

Type compatibility / type equivalence

- Compatibility is the more useful concept, because it tells you what you can DO
- The terms are often used interchangeably

Course Evaluations

- Course evaluations are available on Blackboard
- Be sure to complete all evaluations for all classes
- Only 18% have completed the COMP360 survey



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