

Semantics

COMP360

“Peace cannot be kept by force; it can only be achieved by understanding.”

Albert Einstein

Snowflake Parser

- A recursive descent parser for the Snowflake language is due by noon on Friday, February 17, 2017
- This parser should work with your Snowflake scanner
- The parser must determine if the input Snowflake program has proper syntax
- An appropriate error message (showing line and column) must be displayed if there is a syntax error

Parsing Methods

- Top Down Parsing
 - Traces a leftmost derivation
 - Builds a parse tree in preorder
 - LL and Recursive Descent are top down
- Bottom Up Parsing
 - Builds a parse tree starting at the leaves
 - Produces the reverse of a rightmost derivation
 - LR and LR-K algorithms are bottom up

Recursive Descent

- Relatively easy algorithm to implement
- Functions or methods are written to represent each BNF production
- The goal of the recursive descent parser is to create a parse tree whose leaves (left to right) are the tokens, in order, created by the lexical scanner

Input function

- There is an input function that gets the next token from the lexical scan

```
Token getNextToken() {  
    if (curTok >= tokenList.size()) {  
        return special EOF Token;    // no more tokens  
    }  
    return tokenList.get(curTok++);  
}
```

- I also create a method, peek(), that returns the next token, but does not remove it from the input queue

Example Grammar

$\text{expr} \rightarrow \text{term} \mid \text{term} + \text{expr}$
 $\qquad \qquad \qquad \mid \text{term} - \text{expr}$

$\text{term} \rightarrow \text{factor} \mid \text{factor} * \text{factor}$
 $\qquad \qquad \qquad \mid \text{factor} / \text{factor}$

$\text{factor} \rightarrow \mathbf{\text{variable}} \mid (\text{expr})$

exp function

```
boolean exp() {  
    int saveLocation = curTok;  
    if (!term()) {  
        curTok = saveLocation;  
        return false;  
    }  
    while (peekIs('+') || peekIs('-')) {  
        getNextToken(); // consume + or -  
        if (!term()) {  
            curTok = saveLocation;  
            return false;  
        }  
    }  
    return true;  
}
```

$\text{expr} \rightarrow \text{term} \mid \text{term} + \text{expr}$
 $\mid \text{term} - \text{expr}$

term function

```
boolean term() {  
    int saveLocation = curTok;  
    if (!factor()) {  
        curTok = saveLocation;  
        return false;  
    }  
    while (peekIs('*') || peekIs('/')) {  
        getNextToken();  
        if (!factor()) {  
            curTok = saveLocation;  
            return false;  
        }  
    }  
    return true;  
}
```

term \rightarrow factor | factor * factor
| factor / factor

Write the factor method

factor → **variable** | (expr)

- Assume there is an `isVariable()` method for objects of the `Token` class

Possible Solution

```
boolean term() {                                     factor → variable | ( expr )
    int saveLocation = curTok;
    Token tok = getNextToken();
    if (tok.isVariable()) return true;
    if(tok.value == '(') {
        if (!expr()) {
            curTok = saveLocation;
            return false;
        }
    }
    tok = getNextToken();
    if(tok.value != ')') throw exception;
    return true;
}
```

Handling Errors

- Most compilers attempt to tell the user all syntax errors in their program
- If the program has incorrect syntax in the beginning, it may be difficult for the parser to know how to restart
- Some compilers will move forward to a token that starts the next obvious statement
- The compiler moves up the parse tree to try again

Only the First Error

- Our compiler can stop when it detects an error
- You must report where you think the error is
- An easy way for recursive descent parsers to respond to an error is to throw an exception
- All methods should throw the exception except the main which should catch it and print an error message

java.text.ParseException

- java.text.ParseException is a convenient exception to throw
- The constructor allows a message string and an int
ParseException(String s, int curToken)
- The integer can be an index into the token list indicating where the error was detected
- Alternately, global variables can be set to the error message and location

Errors in Different Phases

- When an error is detected in one phase of a compiler, it will usually stop and not perform any following phases
- If you get a simple error detected by the scanner, it will not run the parser
- Sometimes you will compile a program and get only one error. After you fix that error, you get many more

What made the Scanner program difficult?

- A. I managed to write the program correctly
- B. I don't understand the concept of scanning
- C. I understand scanning, but could not make a program do it
- D. Other factors (i.e. time) prevented me from doing it

Semantic Analysis

- Semantic Analysis checks for non-syntax errors
- Some books define Semantic Analysis as checking all the errors you can't define in BNF

What errors exist in this program?

```
main( ) {  
    short mouse = badProg( 47 );  
}  
  
int badProg( String cat ) {  
    double dog = cat * 1.23;  
    return dog;  
}
```

Some Errors

- The main passes an integer to a String parameter
- The string is multiplied by a double
- The int method returns a double
- The main stores an int result in a short

Which errors would be caught by the parser?

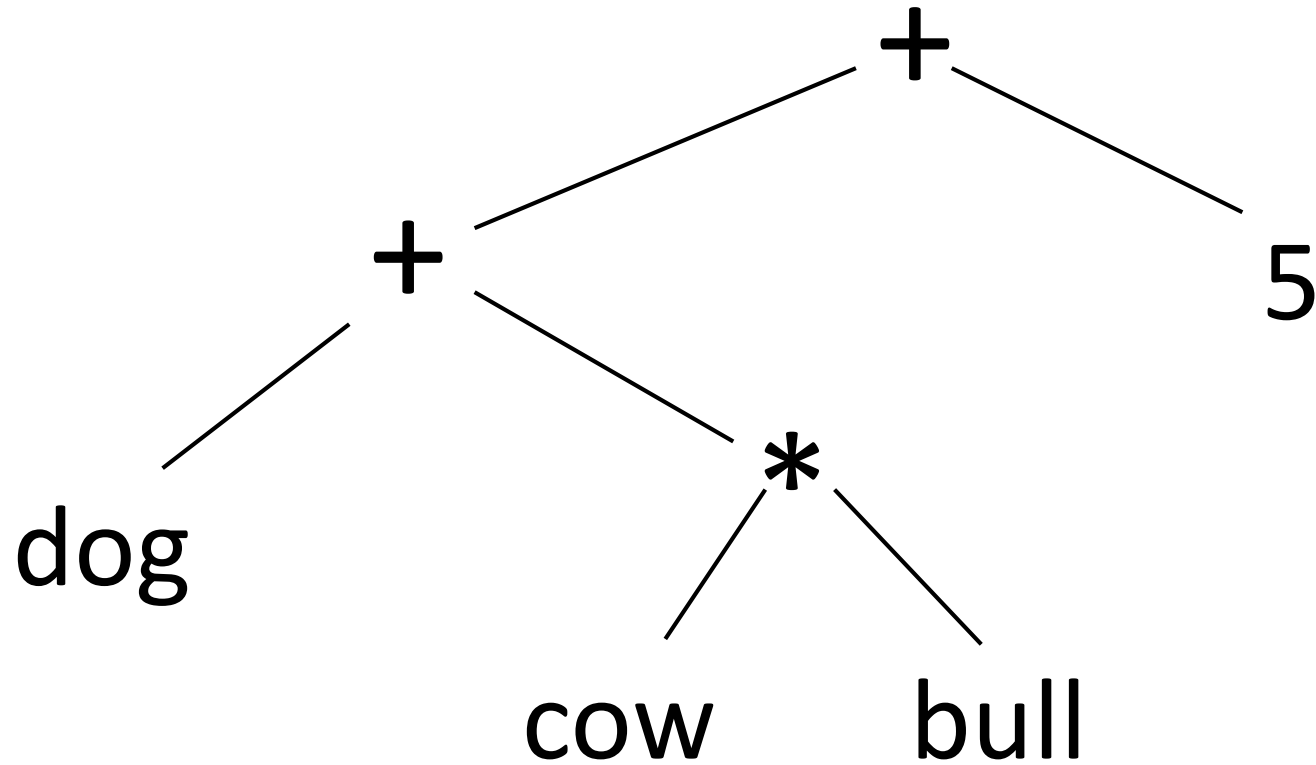
- A. The main passes an integer to a String parameter
- B. The string is multiplied by a double
- C. The int method returns a double
- D. The main stores an int result in a short
- E. None

Creating a Parse Tree

- The compiler has to create a parse tree to direct the code generation
- Traversing the parse tree in a postfix order defines the execution order of the program
- Parse trees are built upon the parse derivation trees with additional information

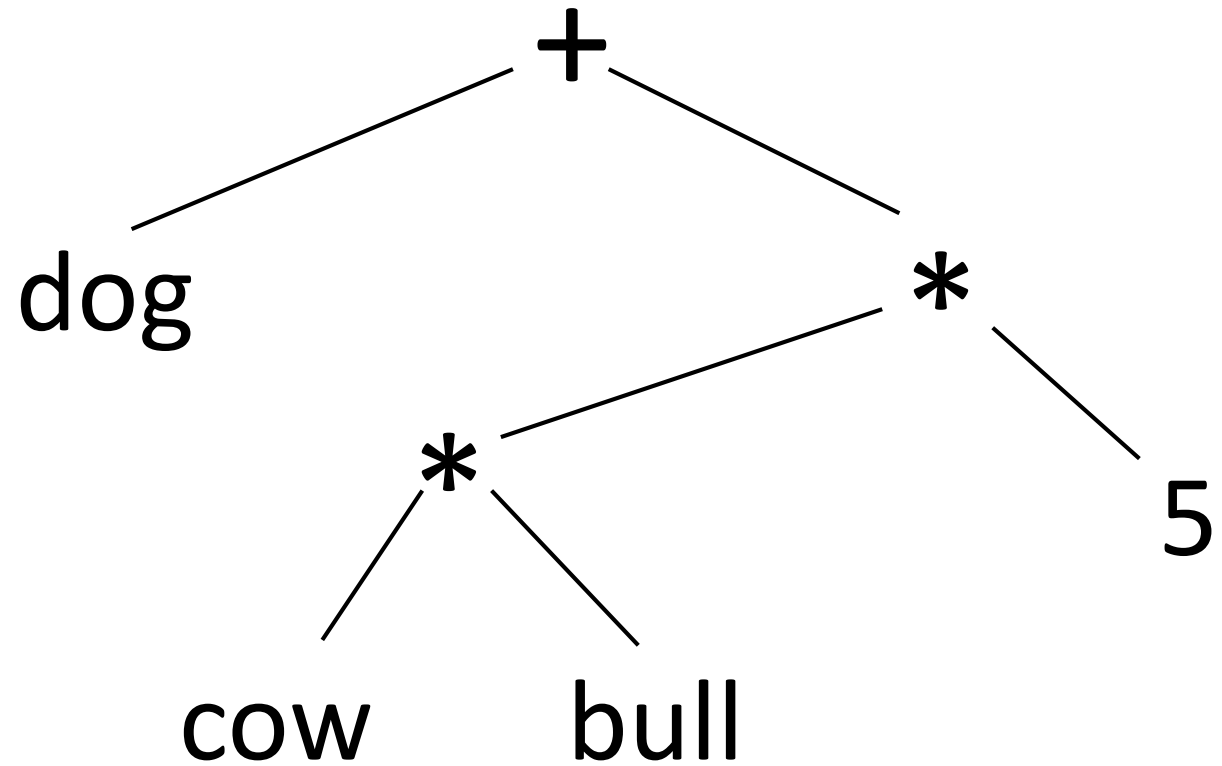
Example Parse Tree

- Consider dog + cow * bull + 5



Example Parse Tree

- Consider $\text{dog} + \text{cow} * \text{bull} * 5$



Boolean Parsing

- The parser you are writing as an assignment only verifies that the input source has correct syntax
- The parser needs to build a parse tree
- Each method of the recursive descent parser needs to add to the parse tree
- Branches are added to the parse tree so the result will execute in the correct order when traversed

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